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# Analyzing Data with Map-Reduce

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## INTRODUCTION TO MAPREDUCE

#### MapReduce Programming Model

- Re-discovered by Google with goals:
	- "Reliability has to come from the software"
	- "How can we make it easy to write distributed programs?"
- A major tool at Google
	- 2.2 million jobs in September 2007 (http://goo.gl/dsnDl)
	- – In 2008 about 100k MapReduce jobs **per day**
		- over 20 petabytes of data processed per day
		- each job occupies about 400 servers

#### Example: Reverse Web-Link Graph



A Problem suitable for MapReduce:

In a set of html-documents, find out which other documents point (via links) to it (for each document)

#### Reverse Web-Link Graph: Solution

c

source

1. For each link to target t found in document source emit (t, source)

2. Sort all pairs by same t's

```
s_{1} _{a} _{a} (a, s_{1})(\mathsf{b},\, \mathsf{s}_1)(\mathtt{C},\,\mathtt{S}_1)\mathbf{s}_{2}\leq \leq \leq(\mathsf{b},\, \mathsf{s}_2)a: (a, s_1)b: (b, s<sub>1</sub>) (b, s<sub>2</sub>) parallel SORT
             c: (c, s_1) (c, s_2)a
                       b
                                                                                            c
                                                                                            b
```
**parallel MAP**

source

3. Results are lists:  $list_t$ (source) for each t

```
list_{a}(s_{1})list<sub>b</sub>(s<sub>1</sub>, s<sub>2</sub>)  parallel REDUCE
list_c(s_2, s_2)
```
#### Distributed MapReduce Frameworks

- • Large data sets require **distribution** 
	- e.g. 1000s of map / reduce tasks in parallel
- Not only Google: **Hadoop** is an open-source implementation



# What Can Go Wrong?

- 1. Master failure
- 2. (Some) workers fail (each with probability p)
	- Case 1: Complete failure
	- –Case 2: Unusual slow execution: **stragglers**
	- Failures "1" are very rare

#### •**Failures 2 are frequent**

- –Workers are commodity (low cost) machines
- – Recall: given n workers, probability of at least one "problem" is

$$
1-(1-p)^n
$$

#### Worker Fault Tolerance in MapReduce

#### **Worker fails completely**

- • Master pings regularly each worker
- • If not responding, worker is marked as "dead"
- •All his (map/red.) jobs inprogress are re-executed

#### **Slow execution: stragglers**

- • Close to completion, masters schedules redundant execution of the remaining *in-progress* jobs
- •Overhead is few % …
- • … but can reduce time-tosolution by up to 1/3

Example: sorting  $10^{10}$  100-byte records (1 TB) with  $\sim$ 1700 workers



#### ONLINE COMPUTING WITH MAPREDUCE

#### Explorative Data Analysis

- Our research at PVS requires a lot of data analysis
- Usually interactive work:
	- 1. Change parameters /code
	- 2. Re-run
	- 3. Evaluate the results
	- 4. Adjust and repeat
- One of the bottlenecks is **time for the mining algorithms to finish**
- => Shortening it reduces the total time of exploration



Orange: Lots of scripting!

#### Time-to-Solution is Important

- Researchers work to a large extend in an "exploratory" way
	- They evaluate last results and decide *then* on the next job(s)
	- – Waiting time for a batch computation to finish is called Time-to-**Solution**
- **Shortening the Time-to-Solution reduces significantly the total time of exploration**



# Online Aggregation

- J. M. Hellerstein, P. Haas and H. Wang introduced in 1997 the concept of Online Aggregation
	- Report online preliminary results (and confidence intervals) for very large queries

SELECT AVG(final\_grade) from grades WHERE course  $name = `CS186`$ GROUP BY major;



- $\bullet$ **Shortens** "Time-to-Decision" in an exploratory data study
	- Allows to cancel a futile query prematurely
	- … Or stop fast if results are precise enough
	- $-$  Helps to identify early how to drill down data

# Incremental-Parallel Data Analysis

- By combining preliminary results with parallelization we get two dimensions of scalability
- We use **Map-Reduce** to combine both approaches



#### **MRStreamer**

- We implemented **MRStreamer** an "enhanced" version of Hadoop
- $\bullet$ It can process data online ("Streaming")
- •… OR in batch mode



#### MRStreamer



#### Features:

- •Batch-mode and incremental (online) processing
- $\bullet$  Efficient shared-memory processing and "flip-aswitch" cluster processing
- $\bullet$ Hadoop-compatible APIs

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### Applications to Data Analysis

#### $\bullet$ Goal: **faster deciding in data analysis studies**

- Estimating whether preliminary solution is stable enough
- $-$  Detecting changes in data profile
- • Example: **online convergence graph**
	- Updates periodically the history of preliminary results



## Algorithms

- Simpler algorithms require **only one MapReduce pass**
	- Aggregation (AVG, SUM, …), Linear regression, PCA, Classification (Naïve Bayes), …
- **Challenging are multi-pass algorithms**
	- For iterative approaches, e.g. clustering via k-means
- $\bullet$ Efficiency dictates changes in algorithms / framework



#### K-Means Clustering Algorithm



# EXTENSIONS AND ALGORITHMS FOR ONLINE MAPREDUCE

#### More Research Problems

- How to enable machine learning algorithms to work **incrementally** (online) and **in parallel**?
- How to help programmers to access / **integrate MapReduce processing in only few lines of code**?
- How to **reduce the inefficiencies** of the frameworks **for smaller data sets**?

#### Framework Inefficiences

- Popular Map-Reduce frameworks like Hadoop are very inefficient for small to medium data sets
	- A job with 5 MB (linear regression) needs on Hadoop 30x longer than on a simple "ad-hoc" MapReduce simulator with 2 threads
	- Hadoop startup time is in the range of 10-20 Seconds
- Key Problem: the code + libraries required for distributed processing introduce overhead not necessary for smaller data sets

#### Efficient Map-Reduce

- • Idea: **one API but resource-specific framework implementations**
	- Dynamic selection of the right infrastructure depending on the input size
- Challenges: coherent APIs and "semantics"
- •First step: MRStreamer
	- –– Both <u>shared-memory</u> and <u>distributed architecture</u>





- • Hadoop is inefficient compared to custom shared-memory code
- Programming in M-R style takes significantly longer
- • Big data sets have specific properties
- Huge data sets are very rare (or inaccessible)

## Which Algorithms Should We Adapt?

- There is a big collection of machine learning alg's running on M-R
	- E.g. **Apache Mahout**
- Adapting them to incremental-parallel processing could be fun …
- But which ones to choose?
- ... Which are really needed?

#### **Mahout currently has**

- Collaborative Filtering
- · User and Item based recommenders
- K-Means, Fuzzy K-Means clustering
- Mean Shift clustering
- Dirichlet process clustering
- **Latent Dirichlet Allocation**
- Singular value decomposition
- Parallel Frequent Pattern mining
- Complementary Naive Bayes classifier
- Random forest decision free based. classifier
- High performance java collections (previously colt collections)
- A vibrant community
- and many more cool stuff to come by this summer thanks to Google summer of code

#### Data Sets: Scale vs. Homogeneity



#### **non-uniform statistical profile / highly structured** 25

#### Consequences for Analyzing Large Data Sets

- Some "small-data" algorithms become useless
	- –They assume a uniform profile in over all data set
- $\bullet$ Other challenges than in "small-data" machine learning
	- –– Identify and recognize rare patterns
	- Split data into sets with homogeneous profile
	- Understand and visualize concept drift
	- Unify models built on parts of the data



# Identifying Concept Drift in Parallel

- $\bullet$  Profile of a modeled phenomenon changes over time
	- –E.g. Profiles of spam emails evolve within days
- •We need a series of models instead of a single one
- •**Concept drift detection** tells us when to switch model
- •Algorithms exist for serial case, what about parallelism?

A joint project with I2R, Singapore



# SEQUENTIAL CONCEPT DRIFT DETECTION

# Classification in a Nutshell

- In **classification** we want to learn from examples model  $f = a$  function from samples (vectors) to elements of a finite set (labels)
- Phase 1: **Training**: we fit/optimize f so that it maps most accurately training samples to their labels
- Phase 2: **Prediction**: f is applied to an unknown sample s to predict its most likely label f(s)



#### Classifiers explained visually

- • If we have only two attributes, we can interpret each sample as a point in  $R^2$
- •Labels are encoded as colors
- • **Training**: finding a suitable subdivision of the plane given the (colored) training points
	- •**Model f** = a compact representation of the subdivision of  $R^2$
- • **Prediction**: given a new sample, find its color = label
- •More metrics  $\Rightarrow$  R<sup>d</sup>



# Incremental Training

#### labeled samples

f learned in order of new training samples

- Assume that we have a long table with labeled samples
- We learn a model f incrementally in order of "new" labeled samples
- How could we detect a concept shift in this scenario?

## Incremental Training with Testing



- After learning f on k first samples, we predict on sample  $s_{k+1}$  and compare prediction f( $s_{k+1}$ ) against true label  $L_{k+1}$
- Then we use labeled  $s_{k+1}$  to further do training of f
- • => So we learn f as before but now also evaluate its accuracy on each new sample (before learning on it)

#### Error Rate

• For each k we can compute the **error rate Err(k)** as:

– $E = \text{Err}(k) = \frac{4}{\pi}$  errors of f until now)/k

• No concept drift: with each new sample f becomes more  $accurate = > Err(k)$  drops





- If the relationship "sample label" changes over time (i.e. we have concept drift), the Err(k) starts to increase after some time!
- => By monitoring Err(k) we can detect concept drift

# Reset after a Concept Drift



- $-$  Start learning a new model (f<sub>2</sub>)
- Reset Err(k)

## Error Rate under Concept Drift



- If Err(k) reaches a critical level:
	- Drop the old model
	- Start learning a new one
	- Reset Err(k)

### Sequential Concept Drift Detection (CDD)

- $\bullet$  The complete algorithm is more complex (link)
- • In addition to the error rate Err(k) we also monitor:
	- –— minimum **P<sub>min</sub>** of Err(k)
	- deviation of Err(k)

We have now two critical levels

- •If  $Err(k) > c_{warm} * (P_{min} + S_{min})$ then **warning signal** is issued
- If  $Err(k) > c_{drift} * (P_{min} + S_{min})$  then <sup>a</sup>**drift signal** is issued



## Sequential CDD /2

- When a warning signal is issued at position **w0**, <sup>a</sup> **"reserve"** classifier **C1** is created
	- – C1 is trained since w0 but not used (yet); the current classifier C0 remains the "main" classifier
- When afterwards a drift signal is issued (at position **dA**):
	- –- We report a concept drift (CD) at d0 and we save or discard C0
	- –C1 replaces C0: C1 becomes the main classifier
	- Minimum P<sub>min</sub> of the error and of variance S<sub>min</sub> are resetted





 $\bullet$  For a distributed version we use also **correctness vectors** (**CV**s) which are sequences of false/true's depending on whether the main classifier was correct or not on a particular sample

– CVs have associated information about their stream position

- $\bullet$ If there are several warnings before a drift, we use the oldest one (not sure: is this also the case in seq. version?)
- $\bullet$  If error(C1) goes below warning level (see **n0** above), we discard a reserve model and create a new (reserve, C2) at next warning (w3)

#### PARALLEL CONCEPT DRIFT DETECTION

# Overall Parallel Architecture /1

• We assume that the incoming stream is multiplexed into **chunks**, each is sent to a mapper / processor



- • Each mapper learns the main (and possibly a reserve) classifier on its chunks (considered as a single stream)
- $\bullet$  Their output (**correctness vector** (**CV**)) contains
	- Correct labels and
	- Labels predicted by the mapper's current classifier



- $\bullet$ Reducer collects the input from all mappers
- •It computes the global error rate Err(k) over all mappers
- $\bullet$ From Err(k) it computes the minima  $P_{min}$ ,  $S_{min}$  and outputs the **levels** normal / warning / drift
	- This is **reducer output stream** (**ROS**)
	- Technicality: The levels are labeled with stream positions



#### Overall Parallel Architecture /3

- $\bullet$ Reducer output stream is "feeded back" to the mappers
- $\bullet$  Each mapper must react to changes in levels of the reducer output:
	- Normal -> warning: start learning reserve classifier
	- Warning -> drift: switch to the reserve classifier



- $\bullet$  Here mappers have sent their outputs to reducer which detected warning at **<sup>w</sup>** and drift at **d**
- $\bullet$ • What should **reducer** do (at d) to mimic the sequential algorithm?
- • Obviously part of the CV after d is useless because it comes from (main) classifiers which should have been replaced at d
- $\bullet$ • Reducer shall *discard the CV after d* and wait for recomputed and "correct" CV (coming from a new classifier)
- $\bullet$  => When reducer receives chunks of the new & correct mapper outputs, it assembles them and continues since d



- $\bullet$  Here mappers have sent their outputs to reducer which detected warning at **<sup>w</sup>** and drift at **d**
- $\bullet$ So upon receiving new input from reducer a mapper does:
	- For warning at position w: starts learning a reserve classifier at w
	- For drift event at d:
		- It switches to the reserve classifier at d
		- It re-computes own output from d to p\_dataEnd and sends it to the reducer

### Summary of Behavior

- Reducer:
	- On drift event at d it discards mapper input after d, resets P<sub>min</sub>, S<sub>min</sub> and waits for re-computed, correct mapper inputs
- $\bullet$  Mapper:
	- On warning event at w: it starts training a reserved classifier since (historical) w
	- On drift event at d: it switches to the reserve classifier, recomputes and re-sends all output to reducer
- Note: drift event at d is like a "sync barrier", it causes all to stop and re-compute everything since d

#### References

- Joos-Hendrik Böse, Artur Andrzejak, Mikael Högqvist: **Beyond Online Aggregation: Parallel and Incremental Data Mining with Online MapReduce**, ACM MDAC 2010, Raleigh, NC, 2010.
- Artur Andrzejak, Joao Bartolo Gomes: **Parallel Concept Drift Detection with Online Map-Reduce**, International Workshop on Knowledge Discovery (KDCloud-2012), at IEEE **ICDM 2012**, Brussels, December 2012.

## THANK YOU. QUESTIONS?