#### RUPRECHT-KARLS-UNIVERSITÄT HEIDELBERG



Prof. Dr. Artur Andrzejak http://pvs.ifi.uni-heidelberg.de/ artur@uni-hd.de

# INTRODUCTION TO MAPREDUCE

#### MapReduce Programming Model

- Re-discovered by Google with goals:
  - "Reliability has to come from the software"
  - "How can we make it easy to write distributed programs?"
- A major tool at Google
  - 2.2 million jobs in September 2007 (http://goo.gl/dsnDl)
  - In 2008 about 100k MapReduce jobs per day
    - over 20 petabytes of data processed per day
    - each job occupies about 400 servers

#### Example: Reverse Web-Link Graph



A Problem suitable for MapReduce:

In a set of html-documents, find out which other documents point (via links) to it (for each document)

#### **Reverse Web-Link Graph: Solution**

source

 For each link to target t found in document source emit (t, source)

2. Sort all pairs by same t's

```
\begin{array}{c|c} s_{1} & (a, s_{1}) \\ c & (b, s_{1}) \\ b & (c, s_{1}) \end{array} & s_{2} & (c, s_{2}) \\ s_{1} & (b, s_{1}) \\ s_{2} & (c, s_{2}) \\ (b, s_{2}) & (b, s_{2}) \\ s_{1} & (b, s_{1}) \\ (b, s_{1}) & (b, s_{2}) \\ s_{2} & (c, s_{2}) \\ (b, s_{2}) & (b, s_{2}) \\ s_{2} & (c, s_{2}) \\ (c, s_{2}) & (c, s_{2}) \end{array}
```

parallel MAP

source

Results are lists:
 list<sub>t</sub>(source) for each t

```
list_{a}(s_{1})
list_{b}(s_{1}, s_{2}) parallel REDUCE
list_{c}(s_{2}, s_{2})
```

#### Distributed MapReduce Frameworks

- Large data sets require **distribution** 
  - e.g. 1000s of map / reduce tasks in parallel
- Not only Google: **Hadoop** is an open-source implementation



# What Can Go Wrong?

- 1. Master failure
- 2. (Some) workers fail (each with probability p)
  - Case 1: Complete failure
  - Case 2: Unusual slow execution: stragglers
- Failures "1" are very rare

#### • Failures 2 are frequent

- Workers are commodity (low cost) machines
- Recall: given n workers, probability of at least one "problem" is

$$1 - (1 - p)^n$$

#### Worker Fault Tolerance in MapReduce

#### Worker fails completely

- Master pings regularly each worker
- If not responding, worker is marked as "dead"
- All his (map/red.) jobs *in- progress* are re-executed

#### Slow execution: stragglers

- Close to completion, masters schedules redundant execution of the remaining *in-progress* jobs
- Overhead is few % ...
- ... but can reduce time-tosolution by up to 1/3

Example: sorting 10<sup>10</sup> 100-byte records (1 TB) with ~1700 workers



#### ONLINE COMPUTING WITH MAPREDUCE

#### Explorative Data Analysis

- Our research at PVS requires a lot of data analysis
- Usually interactive work:
  - 1. Change parameters /code
  - 2. Re-run
  - 3. Evaluate the results
  - 4. Adjust and repeat
- One of the bottlenecks is time for the mining algorithms to finish
- => Shortening it reduces the total time of exploration



Orange: Lots of scripting!

#### Time-to-Solution is Important

- Researchers work to a large extend in an "exploratory" way
  - They evaluate last results and decide *then* on the next job(s)
  - Waiting time for a batch computation to finish is called Time-to-Solution
- Shortening the Time-to-Solution reduces significantly the total time of exploration



# Online Aggregation

- J. M. Hellerstein, P. Haas and H. Wang introduced in 1997 the concept of Online Aggregation
  - Report online preliminary results (and confidence intervals) for very large queries

SELECT AVG(final\_grade) from grades WHERE course\_name = `CS186` GROUP BY major;

🖙 Postgres95 Online Aggregation Interface 📃 🛛 🔀									
	major	AVG	Confidence	Interval					
0	1	2.27216	95	0.160417					
0	2	2.56146	95	0.160417					
0	3	2.66702	95	0.160417					
0	4	2.86235	95	0.160417					
0	5	3.12048	95	0.160417					
0	9	2.89645	95	0.160417					
Cancel All 14% done									

- Shortens "Time-to-Decision" in an exploratory data study
  - Allows to cancel a futile query prematurely
  - ... Or stop fast if results are precise enough
  - Helps to identify early how to drill down data

# Incremental-Parallel Data Analysis

- By combining preliminary results <u>with</u> parallelization we get two dimensions of scalability
- We use Map-Reduce to combine both approaches



#### **MRStreamer**

- We implemented MRStreamer an "enhanced" version of Hadoop
- It can process data <u>online</u> ("Streaming")
- ... OR in batch mode



#### MRStreamer



#### Features:

- Batch-mode and incremental (online) processing
- Efficient shared-memory processing and "flip-aswitch" cluster processing
- Hadoop-compatible APIs

http://pvs.ifi.uni-heidelberg.de/software/mrstreamer/ 15

# Applications to Data Analysis

#### Goal: faster deciding in data analysis studies

- Estimating whether preliminary solution is stable enough
- Detecting changes in data profile
- Example: online convergence graph
  - Updates periodically the history of preliminary results



# Algorithms

- Simpler algorithms require **only** *one* **MapReduce pass** 
  - Aggregation (AVG, SUM, ...), Linear regression, PCA, Classification (Naïve Bayes), ...
- Challenging are multi-pass algorithms
  - For iterative approaches, e.g. clustering via k-means
- Efficiency dictates changes in algorithms / framework



#### **K-Means Clustering Algorithm**



# EXTENSIONS AND ALGORITHMS FOR ONLINE MAPREDUCE

#### More Research Problems

- How to enable machine learning algorithms to work incrementally (online) <u>and</u> in parallel?
- How to help programmers to access / integrate
   MapReduce processing in only few lines of code?
- How to reduce the inefficiencies of the frameworks for smaller data sets?

#### Framework Inefficiences

- Popular Map-Reduce frameworks like Hadoop are very inefficient for small to medium data sets
  - A job with 5 MB (linear regression) needs on Hadoop 30x longer than on a simple "ad-hoc" MapReduce simulator with 2 threads
  - Hadoop startup time is in the range of 10-20 Seconds
- Key Problem: the code + libraries required for distributed processing introduce overhead not necessary for smaller data sets

#### **Efficient Map-Reduce**

- Idea: one API but resource-specific framework implementations
  - Dynamic selection of the right infrastructure depending on the input size
- Challenges: coherent APIs and "semantics"
- First step: <u>MRStreamer</u>
  - Both shared-memory and distributed architecture





- Hadoop is inefficient compared to custom shared-memory code
- Programming in M-R style takes significantly longer
- Big data sets have specific properties
- Huge data sets are very rare (or inaccessible)

#### Which Algorithms Should We Adapt?

- There is a big collection of machine learning alg's running on M-R
  - E.g. Apache Mahout
- Adapting them to incremental-parallel processing could be fun ...
- But which ones to choose?
- ... Which are really needed?

#### Mahout currently has

- Collaborative Filtering
- User and Item based recommenders
- K-Means, Fuzzy K-Means clustering
- Mean Shift clustering
- Dirichlet process clustering
- Latent Dirichlet Allocation
- Singular value decomposition
- Parallel Frequent Pattern mining
- · Complementary Naive Bayes classifier
- Random forest decision tree based classifier
- High performance java collections (previously colt collections)
- A vibrant community
- and many more cool stuff to come by this summer thanks to Google summer of code

## Data Sets: Scale vs. Homogeneity



#### **non**-uniform statistical profile / highly structured<sub>25</sub>

#### Consequences for Analyzing Large Data Sets

- Some "small-data" algorithms become useless
  - They assume a uniform profile in over all data set
- Other challenges than in "small-data" machine learning
  - Identify and recognize rare patterns
  - Split data into sets with homogeneous profile
  - Understand and visualize concept drift
  - Unify models built on parts of the data



# Identifying Concept Drift in Parallel

- Profile of a modeled phenomenon changes over time
  - E.g. Profiles of spam emails evolve within days
- We need a series of models instead of a single one
- Concept drift detection tells us when to switch model
- Algorithms exist for serial case, <u>what about parallelism</u>?

A joint project with I2R, Singapore



# SEQUENTIAL CONCEPT DRIFT DETECTION

# Classification in a Nutshell

- In classification we want to learn from examples model f = a function from samples (vectors) to elements of a finite set (labels)
- Phase 1: **Training**: we fit/optimize f so that it maps most accurately training samples to their labels
- Phase 2: Prediction: f is applied to an unknown sample s to predict its most likely label f(s)

	Attribute 1	•••	Attribute k	Labels	
Training Sample 1	Thursday		1000	ok	h
•••	•••				-1. training
Training Sample N	Sunday		106	defect	
Unknown Sample	Monday		50000	?	← 2. prediction

#### Classifiers explained visually

- If we have only two attributes, we can interpret each sample as <u>a point in R<sup>2</sup></u>
- Labels are encoded as colors
- **Training**: finding a suitable subdivision of the plane given the (colored) training points
  - Model f = a compact representation of the subdivision of R<sup>2</sup>
- Prediction: given a new sample, find its color = label
- More metrics => R<sup>d</sup>



# Incremental Training

#### labeled samples

f learned in order of new training samples

- Assume that we have a long table with labeled samples
- We learn a model f incrementally in order of "new" labeled samples
- How could we detect a concept shift in this scenario?

#### Incremental Training with Testing



- After learning f on k first samples, we predict on sample s<sub>k+1</sub> and compare prediction f(s<sub>k+1</sub>) against true label L<sub>k+1</sub>
- Then we use labeled  $s_{k+1}$  to further do training of f
- => So we learn f as before but now also evaluate its accuracy on each new sample (before learning on it)

#### Error Rate

• For each k we can compute the error rate Err(k) as:

- Err(k) = (# errors of f until now)/k

 No concept drift: with each new sample f becomes more accurate => Err(k) drops





- If the relationship "sample label" changes over time (i.e. we have concept drift), the Err(k) starts to increase after some time!
- => By monitoring Err(k) we can detect concept drift

# Reset after a Concept Drift data collected over time $f_1$ $f_2$ $f_3$ Err(k) We need a <u>new model</u> after a concept drift! • If Err(k) reaches a critical level: - Drop the old model $(f_1)$ - Start learning a new model $(f_2)$

Reset Err(k)

# Error Rate under Concept Drift



- We need a new model after a concept drift!
- If Err(k) reaches a critical level:
  - Drop the old model
  - Start learning a new one
  - Reset Err(k)

#### Sequential Concept Drift Detection (CDD)

- The complete algorithm is • more complex (link)
- In addition to the error rate • Err(k) we also monitor:
  - minimum  $P_{min}$  of Err(k)
  - minimum S<sub>min</sub> of the std. deviation of Err(k)

We have now two critical levels

- If Err(k) >  $c_{warn}$  \* ( $P_{min}$  +  $S_{min}$ ) then warning signal is issued
- If  $Err(k) > C_{drift} * (P_{min} + S_{min})$  then a drift signal is issued



#### Sequential CDD /2

- When a warning signal is issued at position w0, a "reserve" classifier C1 is created
  - C1 is trained since w0 but not used (yet); the current classifier
     C0 remains the "main" classifier
- When afterwards a drift signal is issued (at position dA):
  - We report a concept drift (CD) at d0 and we save or discard C0
  - C1 replaces C0: C1 becomes the main classifier
  - Minimum  $P_{\text{min}}$  of the error and of variance  $S_{\text{min}}$  are resetted





 For a distributed version we use also correctness vectors (CVs) which are sequences of false/true's depending on whether the main classifier was correct or not on a particular sample

- CVs have associated information about their stream position

- If there are several warnings before a drift, we use the <u>oldest</u> one (not sure: is this also the case in seq. version?)
- If error(C1) goes below warning level (see **n0** above), we discard a reserve model and create a new (reserve, C2) at next warning (w3)

#### PARALLEL CONCEPT DRIFT DETECTION

# Overall Parallel Architecture /1

• We assume that the incoming stream is multiplexed into chunks, each is sent to a mapper / processor



- Each mapper learns the main (and possibly a reserve) classifier on its chunks (considered as a single stream)
- Their output (correctness vector (CV)) contains
  - Correct labels and
  - Labels predicted by the mapper's current classifier



- Reducer collects the input from all mappers
- It computes the <u>global</u> error rate Err(k) over all mappers
- From Err(k) it computes the minima P<sub>min</sub>, S<sub>min</sub> and outputs the levels normal / warning / drift
  - This is reducer output stream (ROS)
  - Technicality: The levels are labeled with stream positions



#### Overall Parallel Architecture /3

- Reducer output stream is "feeded back" to the mappers
- Each mapper must react to changes in levels of the reducer output:
  - Normal -> warning: start learning reserve classifier
  - Warning -> drift: switch to the reserve classifier



- Here mappers have sent their outputs to reducer which detected warning at w and drift at d
- What should <u>reducer</u> do (at d) to mimic the sequential algorithm?
- Obviously <u>part of the CV after d is useless</u> because it comes from (main) classifiers which should have been replaced at d
- Reducer shall <u>discard the CV after d</u> and wait for recomputed and "correct" CV (coming from a new classifier)
- => When reducer receives chunks of the new & correct mapper outputs, it assembles them and continues since d



- Here mappers have sent their outputs to reducer which detected warning at w and drift at d
- So upon receiving new input from reducer a <u>mapper</u> does:
  - For warning at position w: starts learning a reserve classifier at w
  - For drift event at d:
    - It switches to the reserve classifier at d
    - It re-computes own output from d to p\_dataEnd and sends it to the reducer

#### Summary of Behavior

- Reducer:
  - On drift event at d it discards mapper input after d, resets
     P<sub>min</sub>, S<sub>min</sub> and waits for re-computed, correct mapper inputs
- Mapper:
  - On warning event at w: it starts training a reserved classifier since (historical) w
  - On drift event at d: it switches to the reserve classifier, recomputes and re-sends all output to reducer
- Note: drift event at d is like a "sync barrier", it causes all to stop and re-compute everything since d

#### References

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- Artur Andrzejak, Joao Bartolo Gomes: *Parallel Concept Drift Detection with Online Map-Reduce*, International Workshop on Knowledge Discovery (KDCloud-2012), at IEEE ICDM 2012, Brussels, December 2012.

# THANK YOU. QUESTIONS?